



# Advanced Data Structures and Algorithms

Associate Professor Dr. Raed Ibraheem Hamed

Computer Science Department  
College of Science and Technology  
University of Human Development,

2015 – 2016

Department of Computer Science \_ UHD

# What this Lecture is about:

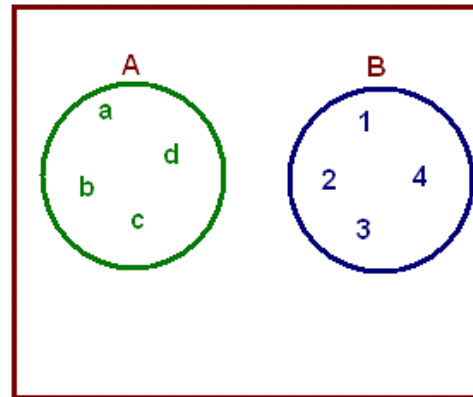
---

- A disjoint-set data structure
- Disjoint Set Operations
- An Application of Disjoint-Set
- An Application of Disjoint-Set Data Structures
- Linked-List Implementation
- Linked-lists for two sets
- Disjoint-set Implementation: Forests
- Algorithm for Disjoint-Set
- Example

# A disjoint-set data structure

---

- A disjoint-set is a collection  $\mathbf{S}=\{S_1, S_2, \dots, S_k\}$  of distinct dynamic sets.
- Is a data structure that keeps track of a set of elements partitioned into a number  $f$  disjoint subsets.
- Each set is identified by a member of the set, called *representative*.



Two disjoint sets A and B

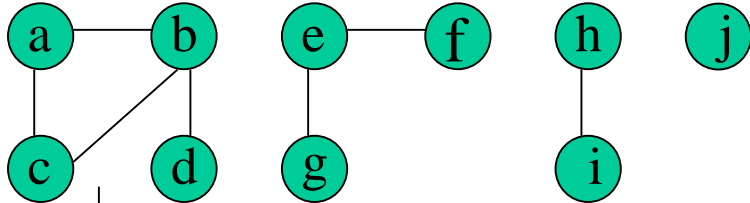
# Disjoint set operations:

---

- **MAKE-SET( $x$ )**: create a new set with only  $x$ . assume  $x$  is not already in some other set.
- **UNION( $x,y$ )**: combine the two sets containing  $x$  and  $y$  into one new set. A new representative is selected.
- **FIND-SET( $x$ )**: returns a pointer to the representative of the unique set containing  $x$ . *Find: Determine which subset a particular element is in.*

# An Application of Disjoint-Set Data Structures

*Determining the connected components of an undirected graph  $G=(V,E)$*

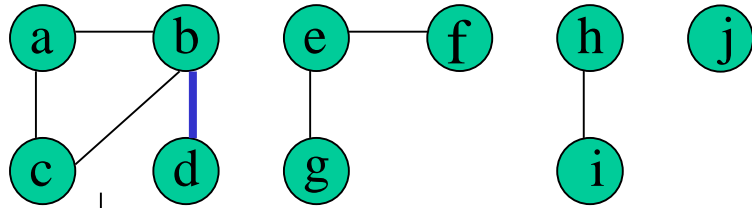


---

Initial	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
---------	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

# An Application of Disjoint-Set Data Structures

*Determining the connected components of an undirected graph  $G=(V,E)$*

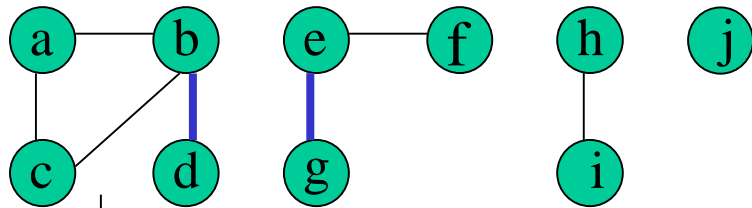


---

Initial	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
(b, d)	{a}	{b, d}	{c}		{e}	{f}	{g}	{h}	{i}	{j}

# An Application of Disjoint-Set Data Structures

*Determining the connected components of an undirected graph  $G=(V,E)$*

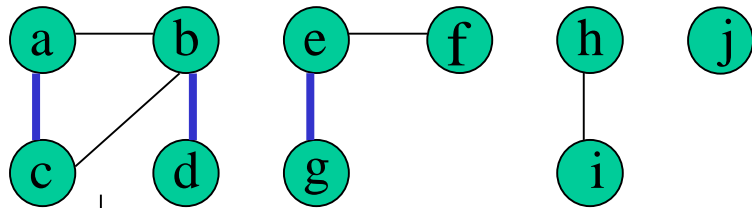


---

Initial	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
(b, d)	{a}	{b, d}	{c}		{e}	{f}	{g}	{h}	{i}	{j}
(e, g)	{a}	{b, d}	{c}		{e, g}	{f}		{h}	{i}	{j}

# An Application of Disjoint-Set Data Structures

*Determining the connected components of an undirected graph  $G=(V,E)$*

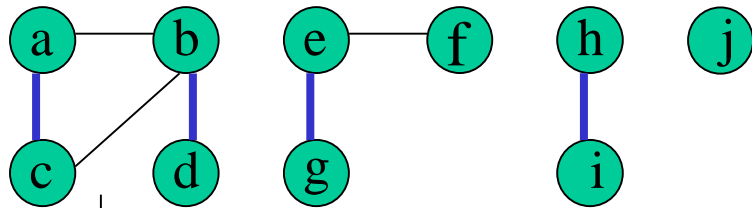


Initial	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
(b, d)	{a}	{b, d}	{c}		{e}	{f}	{g}	{h}	{i}	{j}
(e, g)	{a}	{b, d}	{c}		{e, g}	{f}		{h}	{i}	{j}
(a, c)	{a, c}	{b, d}			{e, g}	{f}		{h}	{i}	{j}



# An Application of Disjoint-Set Data Structures

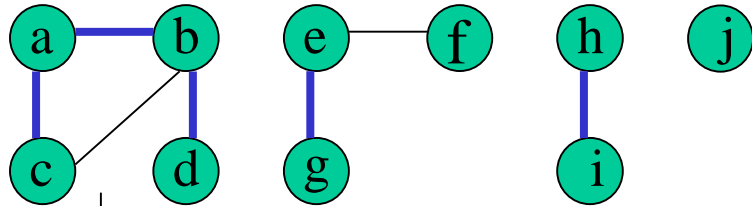
*Determining the connected components of an undirected graph  $G=(V,E)$*



Initial	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
(b, d)	{a}	{b, d}	{c}		{e}	{f}	{g}	{h}	{i}	{j}
(e, g)	{a}	{b, d}	{c}		{e, g}	{f}		{h}	{i}	{j}
(a, c)	{a, c}	{b, d}			{e, g}	{f}		{h}	{i}	{j}
(h, i)	{a, c}	{b, d}			{e, g}	{f}		{h, i}		{j}

# An Application of Disjoint-Set Data Structures

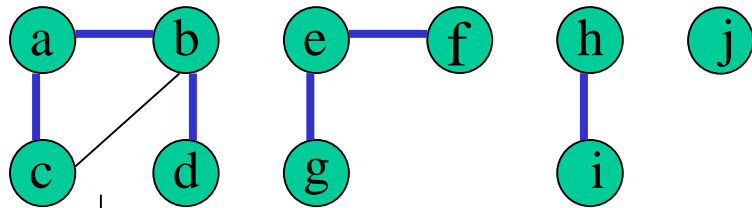
*Determining the connected components of an undirected graph  $G=(V,E)$*



Initial	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
(b, d)	{a}	{b, d}	{c}		{e}	{f}	{g}	{h}	{i}	{j}
(e, g)	{a}	{b, d}	{c}		{e, g}	{f}		{h}	{i}	{j}
(a, c)	{a, c}	{b, d}			{e, g}	{f}		{h}	{i}	{j}
(h, i)	{a, c}	{b, d}			{e, g}	{f}		{h, i}		{j}
(a, b)	{a, b, c, d}				{e, g}	{f}		{h, i}		{j}

# An Application of Disjoint-Set Data Structures

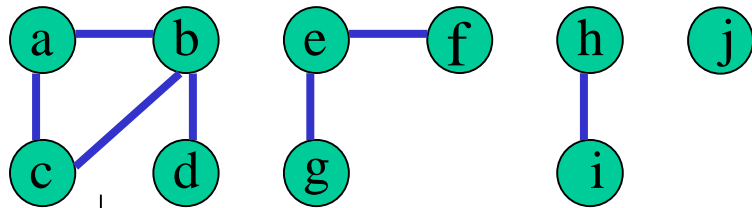
*Determining the connected components of an undirected graph  $G=(V,E)$*



Initial	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
(b, d)	{a}	{b, d}	{c}		{e}	{f}	{g}	{h}	{i}	{j}
(e, g)	{a}	{b, d}	{c}		{e, g}	{f}		{h}	{i}	{j}
(a, c)	{a, c}	{b, d}			{e, g}	{f}		{h}	{i}	{j}
(h, i)	{a, c}	{b, d}			{e, g}	{f}		{h, i}		{j}
(a, b)	{a, b, c, d}				{e, g}	{f}		{h, i}		{j}
(e, f)	{a, b, c, d}				{e, f, g}			{h, i}		{j}

# An Application of Disjoint-Set Data Structures

*Determining the connected components of an undirected graph  $G=(V,E)$*



Initial	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
(b, d)	{a}	{b, d}	{c}		{e}	{f}	{g}	{h}	{i}	{j}
(e, g)	{a}	{b, d}	{c}		{e, g}	{f}		{h}	{i}	{j}
(a, c)	{a, c}	{b, d}			{e, g}	{f}		{h}	{i}	{j}
(h, i)	{a, c}	{b, d}			{e, g}	{f}		{h, i}		{j}
(a, b)	{a, b, c, d}				{e, g}	{f}		{h, i}		{j}
(e, f)	{a, b, c, d}				{e, f, g}			{h, i}		{j}
(b, c)	{a, b, c, d}				{e, f, g}			{h, i}		{j}

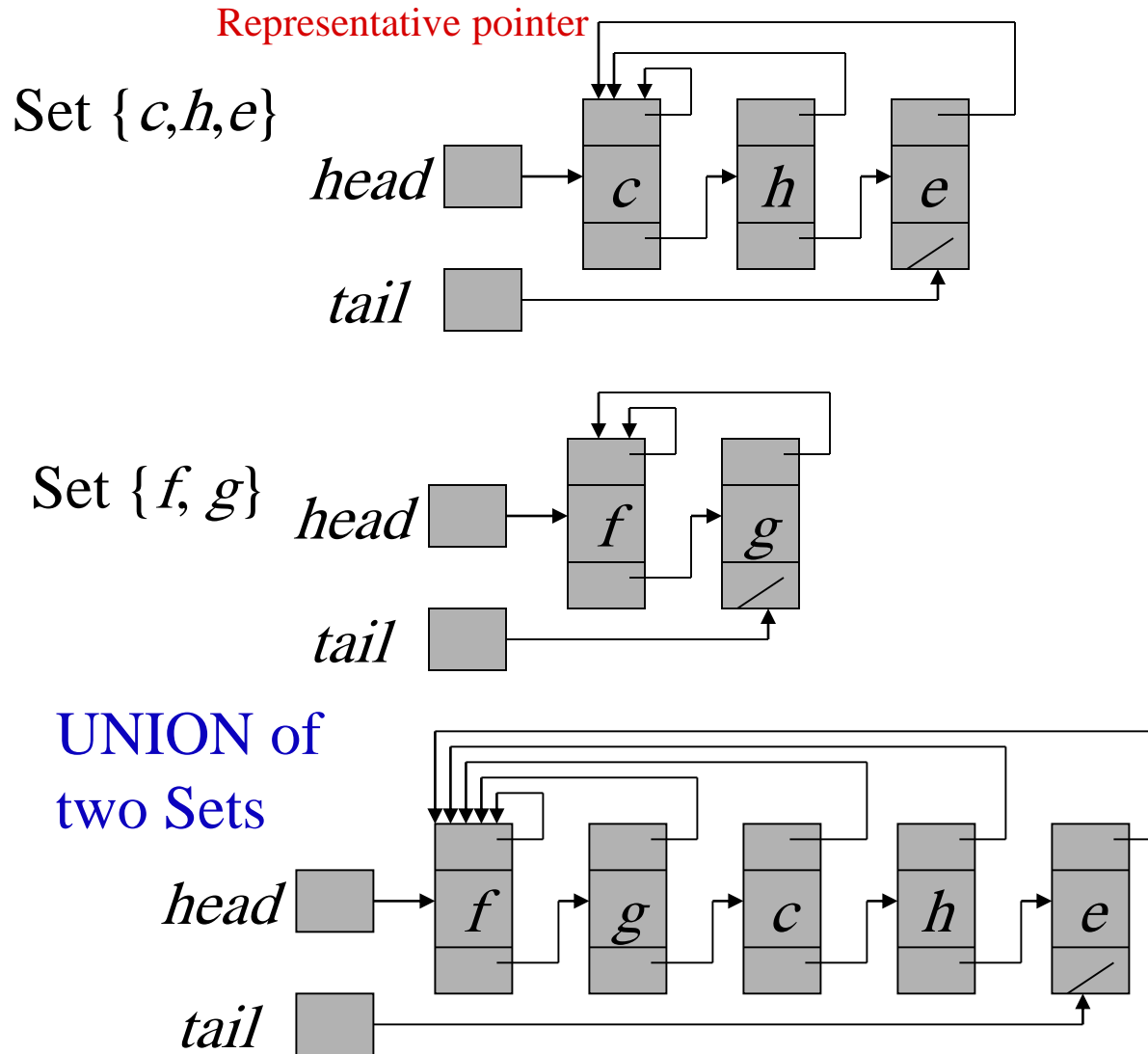
# Linked-List Implementation

---

- Each set as a linked-list, with head and tail, and each node contains value, next node pointer and back-to-representative pointer.
- *Example:*
- MAKE-SET costs  $O(1)$ : just create a single element list.
- FIND-SET costs  $O(1)$ : just return back-to-representative pointer.

# Linked-lists for two sets

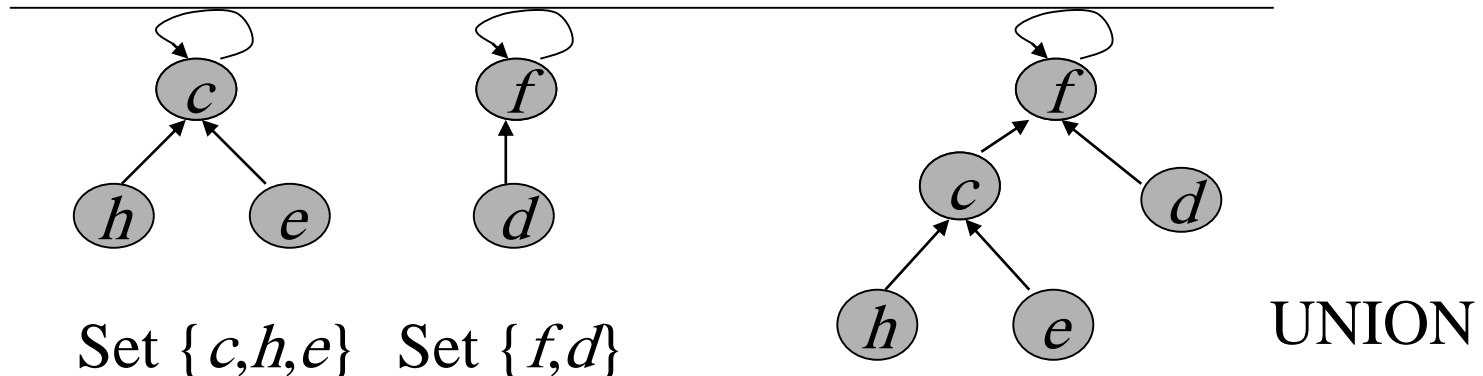
---



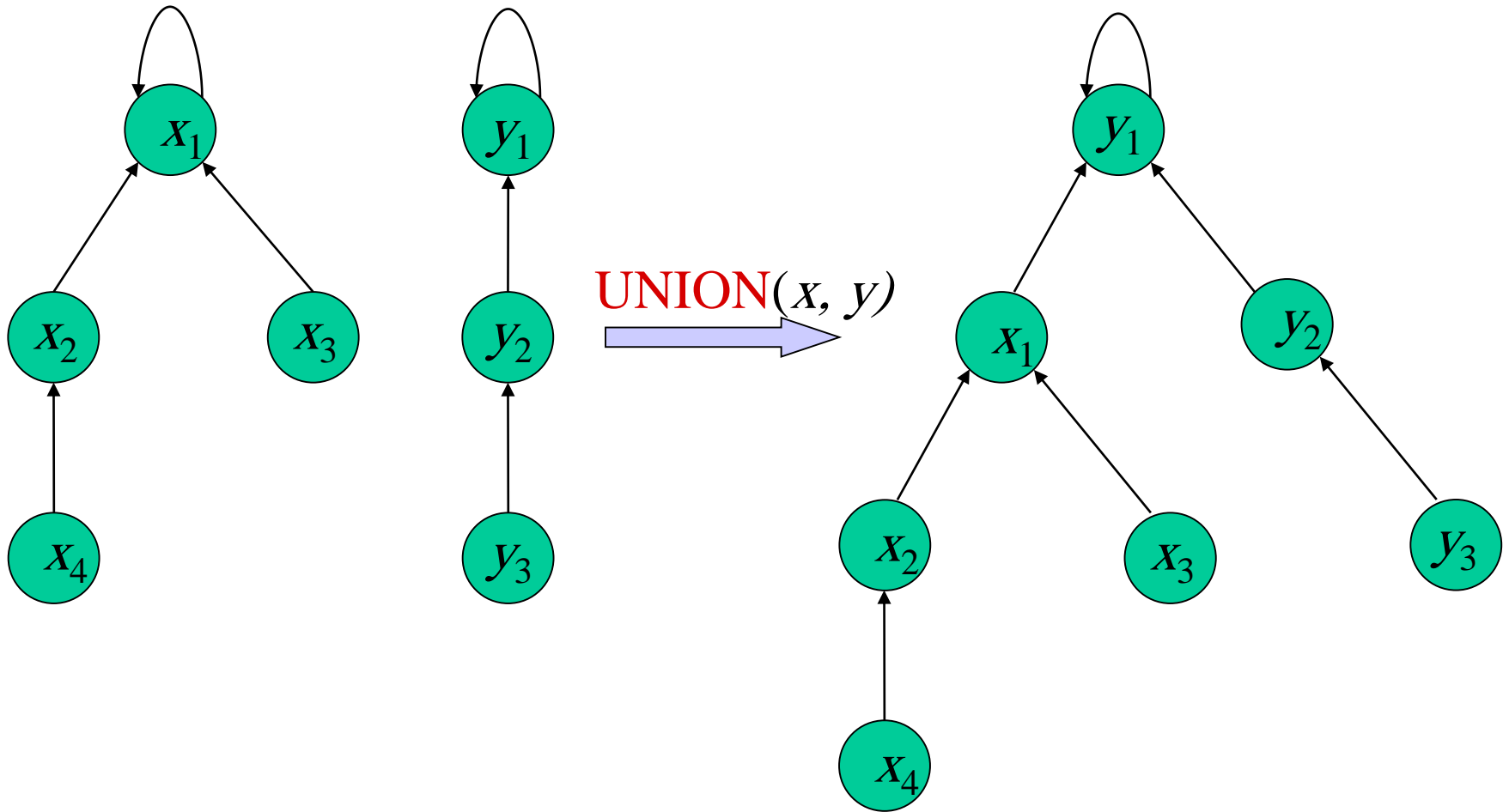
# Disjoint-set Implementation: Forests

---

- Rooted trees, each tree is a set, root is the **representative**. Each node points to its parent. Root points to itself.



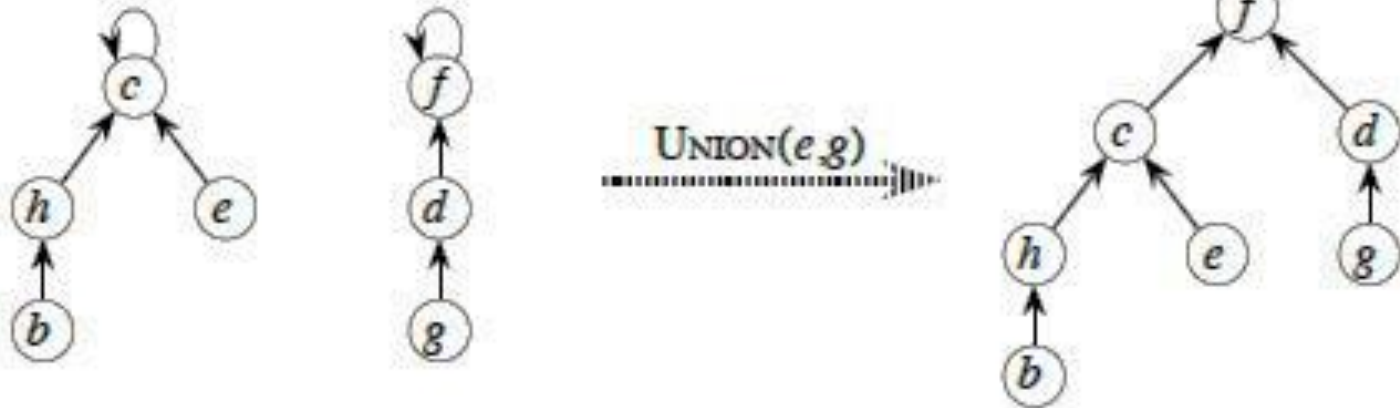
# Disjoint Set Forests





# Disjoint Set Forests

---



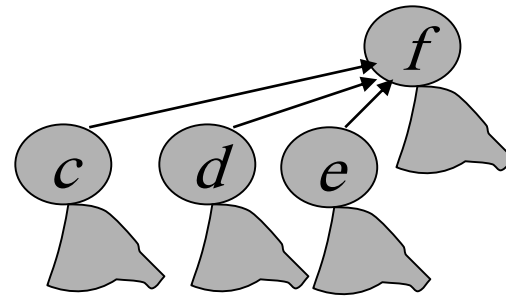
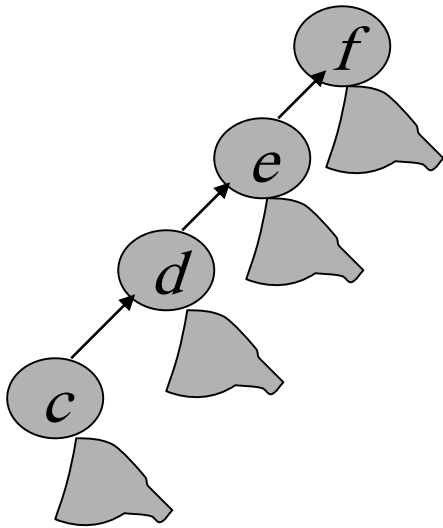
# Path Compression

---

**Path Compression:** used in  $\text{FIND-SET}(x)$  operation, make each node in the path from  $x$  to the root directly point to the root. Thus reduce the tree height.

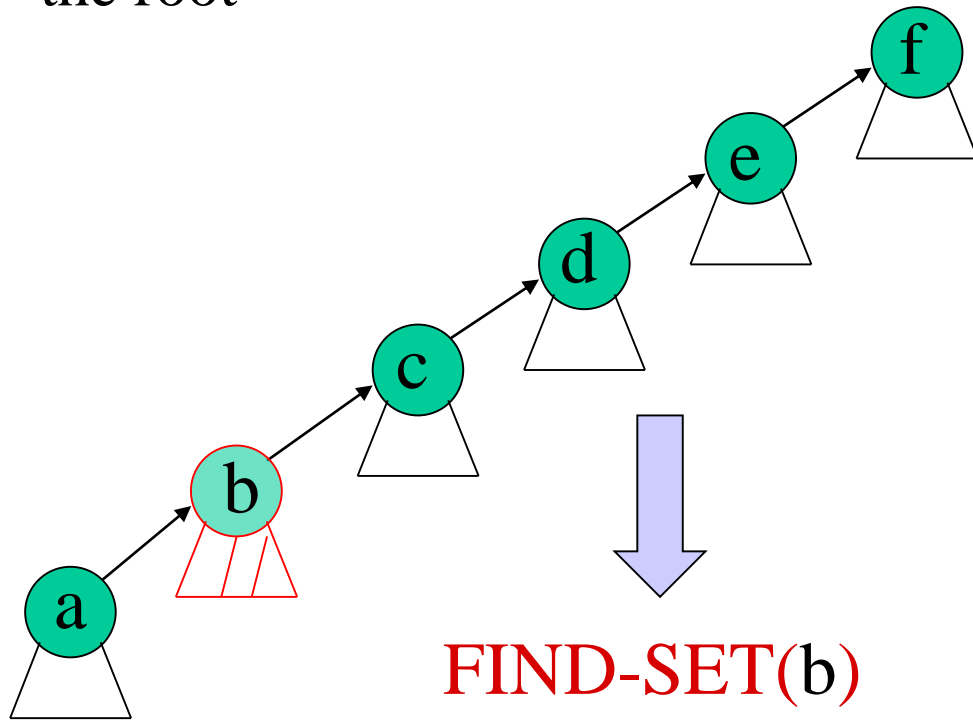
# Path Compression

---



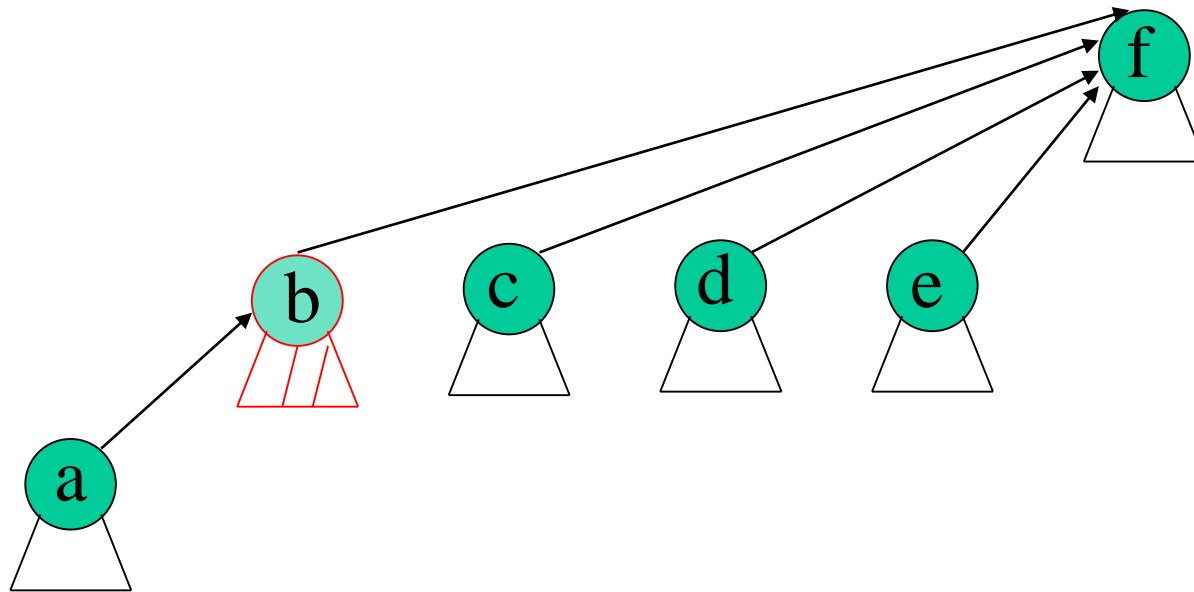
# Path Compression

- Use it during the **FIND-SET** operations
- Make each node on the **FIND-PATH** to point directly to the root



# Path Compression

Path Compression During **FIND-SET(b)** Operation



# Algorithm for Disjoint-Set

---

```
function MakeSet(x)
```

```
    x.parent := x
```

```
function Find(x)
```

```
    if x.parent == x
```

```
        return x
```

```
    Else
```

```
        return Find(x.parent)
```

```
function Union(x, y)
```

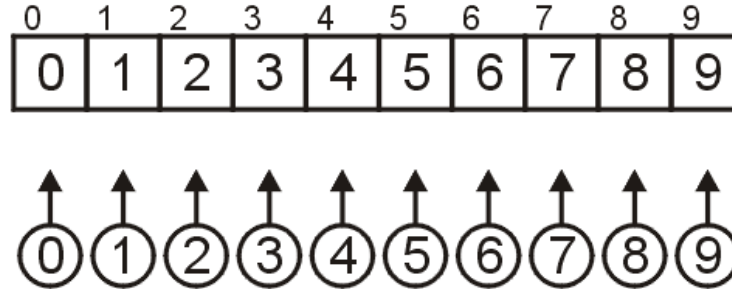
```
    xRoot := Find(x)
```

```
    yRoot := Find(y)
```

```
    xRoot.parent := yRoot
```

# Example

Consider the following disjoint set on the ten decimal digits:



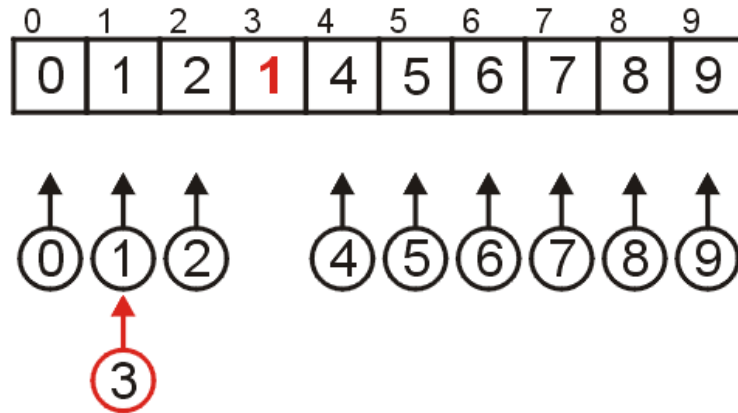
$\{0\}, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\}$

# Example

If we take the union of the sets containing 1 and 3

```
set_union(1, 3);
```

we perform a find on both entries and update the second

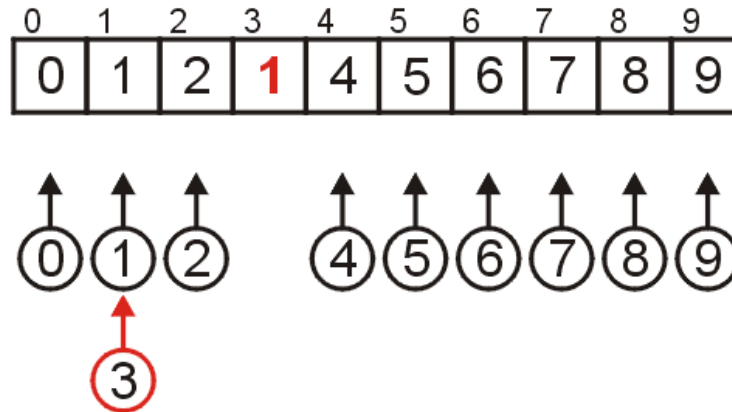


$\{0\}, \{1, 3\}, \{2\}, \{4\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\}$



# Example

Now, `find(1)` and `find(3)` will both return the integer 1

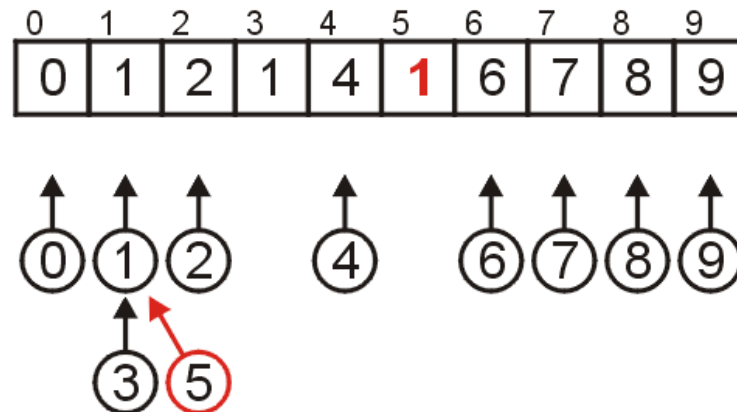


{0}, {1, 3}, {2}, {4}, {5}, {6}, {7}, {8}, {9}

# Example

Next, take the union of the sets containing 3 and 5,  
`set_union(3, 5);`

we perform a find on both entries and update the second



$\{0\}, \{1, 3, 5\}, \{2\}, \{4\}, \{6\}, \{7\}, \{8\}, \{9\}$

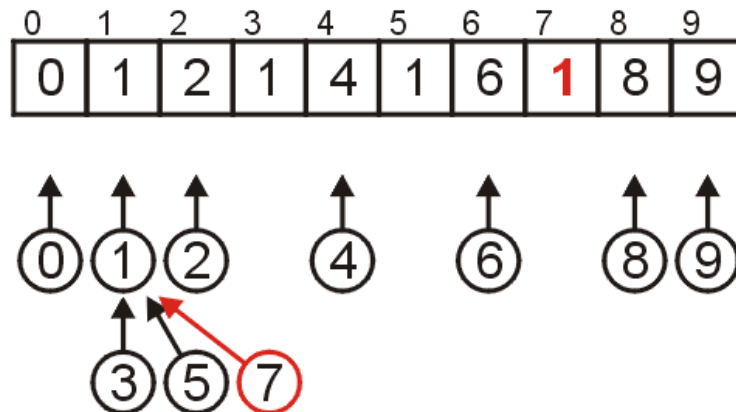
# Example

---

Now, if we take the union of the sets containing 5 and 7

```
set_union(5, 7);
```

we update the value stored in `find(7)` with the value `find(5)`:

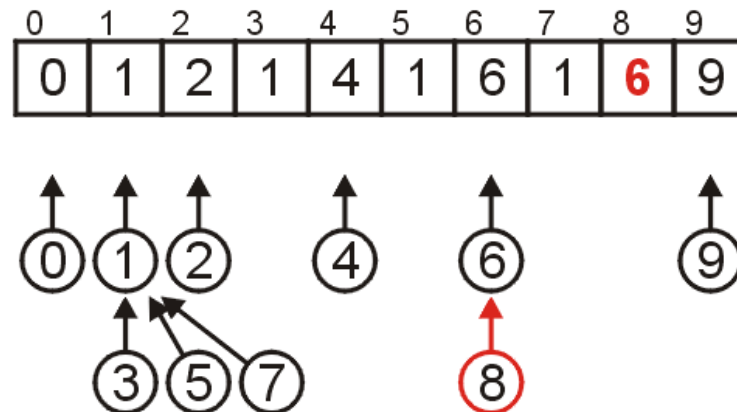


$\{0\}, \{1, 3, 5, 7\}, \{2\}, \{4\}, \{6\}, \{8\}, \{9\}$

# Example

Taking the union of the sets containing 6 and 8  
`set_union(6, 8);`

we update the value stored in `find(8)` with the value  
`find(6)`:



$\{0\}, \{1, 3, 5, 7\}, \{2\}, \{4\}, \{6, 8\}, \{9\}$

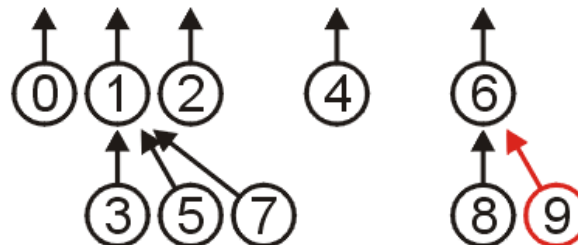
# Example

Taking the union of the sets containing 8 and 9

```
set_union(8, 9);
```

we update the value stored in `find(8)` with the value `find(9)`:

0	1	2	3	4	5	6	7	8	9
0	1	2	1	4	1	6	1	6	6



$\{0\}, \{1, 3, 5, 7\}, \{2\}, \{4\}, \{6, 8, 9\}$

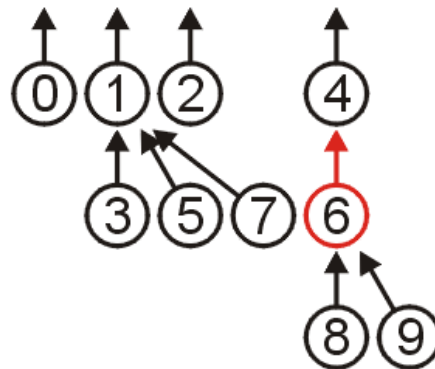
# Example

Taking the union of the sets containing 4 and 8

```
set_union(4, 8);
```

we update the value stored in `find(8)` with the value `find(4)`:

0	1	2	3	4	5	6	7	8	9
0	1	2	1	4	1	4	1	6	6



$\{0\}, \{1, 3, 5, 7\}, \{2\}, \{4, 6, 8, 9\}$

# Example

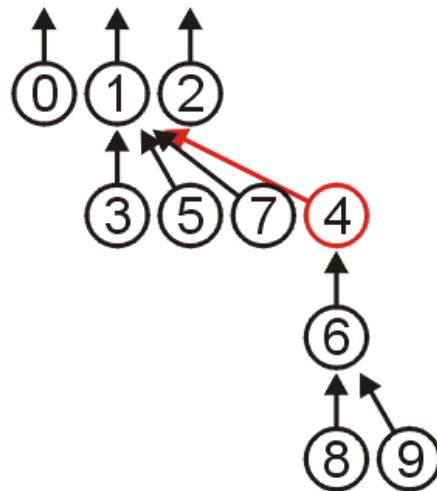
---

Finally, if we take the union of the sets containing 5 and 6

```
set_union(5, 6);
```

we update the entry of `find(6)` with the value of `find(5)`:

0	1	2	3	4	5	6	7	8	9
0	1	2	1	1	1	4	1	6	6



$\{0\}, \{1, 3, 4, 5, 6, 7, 8, 9\}, \{2\}$

---

THANK YOU

???